

# Applicative invariants for Lustre

Amos Robinson, international man of leisure



Chestnut-backed chickadee

# Synchronous languages at an autonomous vehicle startup



# Happy days

- Implemented controllers, drivers, state machines in Lustre
- Proved some reassuring properties
- Ran them in the car
- Worked first time (more or less)

# Success with Kind2

Kind2 is generally great:

- Solid for small programs
- Actively developed
- Friendly, helpful developers



# Sadness with Kind2

But as our programs got bigger...

- Hard to predict whether true properties can be proved
- Modifications difficult, could require serious restructuring to keep proofs
- Spurious failures in CI make us look bad

# What we couldn't prove reliably



Willie wagtail vs grey butcherbird

```
node lastn(const n: int; pred: bool)
returns (out: bool)
let
  count = if pred then (0 -> pre count) + 1 else 0;
  out   = count >= n;
tel
```

```
node lastn(const n: int; pred: bool)
returns (out: bool)
let
  count = if pred then (0 -> pre count) + 1 else 0;
  out   = count >= n;
tel
```

```
function delta_valid(input1: int, input2: int)
returns (ok: bool)
let
  ok = abs(input1 - input2) <= DELTA_MAX
tel
```

```
node signal_valid(input: int)
returns (ok: bool)
(*@contract
guarantee
  not lastn(10, delta_valid(input, 0 -> pre input)) =>
  not ok;
  ...
*)
```

```
node signal_valid(input: int)
returns (ok: bool)
(*@contract
  guarantee
    not lastn(10, delta_valid(input, 0 -> pre input)) =>
    not ok;
  ...
*)
```

```
type SAMPLE      = { adc: int; ... }
const SAMPLE_ZERO = { adc = 0; ... }
```

```
node main(sample: SAMPLE)
returns (engaged: bool; ...)
(*@contract
  guarantee
    not lastn(10, delta_valid(sample.adc, (SAMPLE_ZERO -> pre sample).adc)) =>
    not engaged;
*)
let
  engaged = signal_valid(sample.adc);
  ...
tel
```

# Pen-and-paper proof



```
assume
```

```
  signal_valid.guarantee[input := sample.adc, ok := engaged] :  
    not lastn(10, delta_valid(sample.adc, 0 -> pre sample.adc)) =>  
    not engaged;
```

```
engaged = signal_valid(sample.adc);
```

```
assume
```

```
  signal_valid.guarantee[input := sample.adc, ok := engaged] :  
    not lastn(10, delta_valid(sample.adc, 0 -> pre sample.adc)) =>  
    not engaged;
```

```
engaged = signal_valid(sample.adc);
```

```
SAMPLE_ZERO = { adc = 0; ... };
```

```
assume
```

```
  signal_valid.guarantee[input := sample.adc, ok := engaged] :  
    not lastn(10, delta_valid(sample.adc, 0 -> pre sample.adc)) =>  
    not engaged;
```

```
engaged = signal_valid(sample.adc);
```

```
SAMPLE_ZERO = { adc = 0; ... };
```

---

```
show main.guarantee:
```

```
  not lastn(10, delta_valid(sample.adc, (SAMPLE_ZERO -> pre sample).adc)) =>  
  not engaged;
```

```
assume
```

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  signal_valid.guarantee[input := sample.adc, ok := engaged] :  
    not lastn(10, delta_valid(sample.adc, 0 -> pre sample.adc)) =>  
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```
engaged = signal_valid(sample.adc);
```

```
SAMPLE_ZERO = { adc = 0; ... };
```

---

```
show main.guarantee:
```

```
  not lastn(10, delta_valid(sample.adc, (SAMPLE_ZERO -> pre sample).adc)) =>  
  not engaged;
```

```
rewrite prim_distributes_arrow: forall stream e e', pure function f.
```

```
  f (e -> e') = (f e) -> (f e')
```

```
assume
```

```
  signal_valid.guarantee[input := sample.adc, ok := engaged] :  
    not lastn(10, delta_valid(sample.adc, 0 -> pre sample.adc)) =>  
    not engaged;
```

```
engaged = signal_valid(sample.adc);
```

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SAMPLE_ZERO = { adc = 0; ... };
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```
show main.guarantee:
```

```
  not lastn(10, delta_valid(sample.adc, SAMPLE_ZERO.adc -> (pre sample).adc)) =>  
  not engaged;
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rewrite prim_distributes_arrow: forall stream e e', pure function f.
```

```
  f (e -> e') = (f e) -> (f e')
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```
assume
```

```
  signal_valid.guarantee[input := sample.adc, ok := engaged] :  
    not lastn(10, delta_valid(sample.adc, 0 -> pre sample.adc)) =>  
    not engaged;
```

```
engaged = signal_valid(sample.adc);
```

```
SAMPLE_ZERO = { adc = 0; ... };
```

---

```
show main.guarantee:
```

```
  not lastn(10, delta_valid(sample.adc, SAMPLE_ZERO.adc -> (pre sample).adc)) =>  
  not engaged;
```

```
rewrite prim_distributes_pre: forall stream e, pure function f.
```

```
  f (pre e) = pre (f e)
```

```
assume
```

```
  signal_valid.guarantee[input := sample.adc, ok := engaged] :  
    not lastn(10, delta_valid(sample.adc, 0 -> pre sample.adc)) =>  
    not engaged;
```

```
engaged = signal_valid(sample.adc);
```

```
SAMPLE_ZERO = { adc = 0; ... };
```

---

```
show main.guarantee:
```

```
  not lastn(10, delta_valid(sample.adc, SAMPLE_ZERO.adc -> pre sample.adc)) =>  
  not engaged;
```

```
rewrite prim_distributes_pre: forall stream e, pure function f.
```

```
  f (pre e) = pre (f e)
```

```
assume
```

```
  signal_valid.guarantee[input := sample.adc, ok := engaged] :  
    not lastn(10, delta_valid(sample.adc, 0 -> pre sample.adc)) =>  
    not engaged;
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```
engaged = signal_valid(sample.adc);
```

```
SAMPLE_ZERO = { adc = 0; ... };
```

---

```
show main.guarantee:
```

```
  not lastn(10, delta_valid(sample.adc, 0 -> pre sample.adc)) =>  
  not engaged;
```

```
unfold SAMPLE_ZERO, constant propagation, etc
```

# Transition system proof



Australasian grebes

```
node delta_valid_sample(sample: SAMPLE)
returns (delta_main: bool, delta_signal: bool)
let
    delta_main    = delta_valid(sample.adc, (SAMPLE_ZERO -> pre sample).adc);
    delta_signal = delta_valid(sample.adc, 0 -> pre sample.adc);
tel
```

```
node delta_valid_sample(sample: SAMPLE)
returns (delta_main: bool, delta_signal: bool)
let
  delta_main = delta_valid(sample.adc, (SAMPLE_ZERO -> pre sample).adc);
  delta_signal = delta_valid(sample.adc, 0 -> pre sample.adc);
tel
```

==>

```
lts delta_valid_sample:
type state = { init:          bool;
               delay_main:    SAMPLE;
               delay_signal: int; };

init(s: state): bool =
  s.init == true;

step(s: state, sample: SAMPLE): (state * bool * bool) =
  let
    delay_main           = if s.init then SAMPLE_ZERO else s.delay_main
    delta_main           = delta_valid(sample.adc, delay_main.adc)

    delay_signal         = if s.init then 0 else s.delay_signal
    delta_signal         = delta_valid(sample.adc, delay_signal)

  in
    ({ init          = false;
       delay_main    = sample;
       delay_signal  = sample.adc },
     delta_main, delta_signal)
```

```
node delta_valid_sample(sample: SAMPLE)
returns (delta_main: bool, delta_signal: bool)
let
  delta_main = delta_valid(sample.adc, (SAMPLE_ZERO -> pre sample).adc);
  delta_signal = delta_valid(sample.adc, 0 -> pre sample.adc);
tel
```

==>

```
lts delta_valid_sample:
  type state = { init:          bool;
                 delay_main:    SAMPLE;
                 delay_signal: int; };

  init(s: state): bool =
    s.init == true;

  step(s: state, sample: SAMPLE): (state * bool * bool) =
    let
      delay_main           = if s.init then SAMPLE_ZERO else s.delay_main
      delta_main           = delta_valid(sample.adc, delay_main.adc)

      delay_signal         = if s.init then 0 else s.delay_signal
      delta_signal         = delta_valid(sample.adc, delay_signal)

    in
      ({ init           = false;
         delay_main     = sample;
         delay_signal   = sample.adc },
       delta_main, delta_signal)
```

```
node delta_valid_last2(delta_ok: bool)
returns (last_main: bool, last_signal: bool)
let
    last_main = lastn(10, delta_ok);
    last_signal = lastn(10, delta_ok);
```

```
tel
```

```
node delta_valid_last2(delta_ok: bool)
returns (last_main: bool, last_signal: bool)
let
  last_main = lastn(10, delta_ok);
  last_signal = lastn(10, delta_ok);
```

```
tel
```

```
==>
```

```
lts delta_valid_last2:
  type state = { last_main: lastn.state;
                  last_signal: lastn.state; };

  init(s: state): bool =
    lastn.init(s.last_main) and lastn.init(s.last_signal);

  step(s: state, delta_ok: bool): (state * bool * bool) =
    let
      (state_last_main', last_main) = lastn.step(s.last_main, 10, delta_ok)
      (state_last_signal', last_signal) = lastn.step(s.last_signal, 10, delta_ok)
    in
      ({ init      = false;
         last_main = state_last_main';
         last_signal = state_last_signal' }, last_main, last_signal)
```

```
node delta_valid_last2(delta_ok: bool)
returns (last_main: bool, last_signal: bool)
let
  last_main = lastn(10, delta_ok);
  last_signal = lastn(10, delta_ok);
```

```
tel
```

```
==>
```

```
lts delta_valid_last2:
  type state = { last_main: lastn.state;
                  last_signal: lastn.state; };

  init(s: state): bool =
    lastn.init(s.last_main) and lastn.init(s.last_signal);

  step(s: state, delta_ok: bool): (state * bool * bool) =
    let
      (state_last_main', last_main) = lastn.step(s.last_main, 10, delta_ok)
      (state_last_signal', last_signal) = lastn.step(s.last_signal, 10, delta_ok)
    in
      ({ init      = false;
         last_main = state_last_main';
         last_signal = state_last_signal' }, last_main, last_signal)
--%PROPERTY state_last_main' == state_last_signal'
```

```
node delta_valid_last2(delta_ok: bool)
returns (last_main: bool, last_signal: bool)
let
    last_main = lastn(10, delta_ok);
    last_signal = lastn(10, delta_ok);
    --%PROPERTY ??? = ???;
tel
```

==>

```
lts delta_valid_last2:
  type state = { last_main: lastn.state;
                  last_signal: lastn.state; };

  init(s: state): bool =
    lastn.init(s.last_main) and lastn.init(s.last_signal);

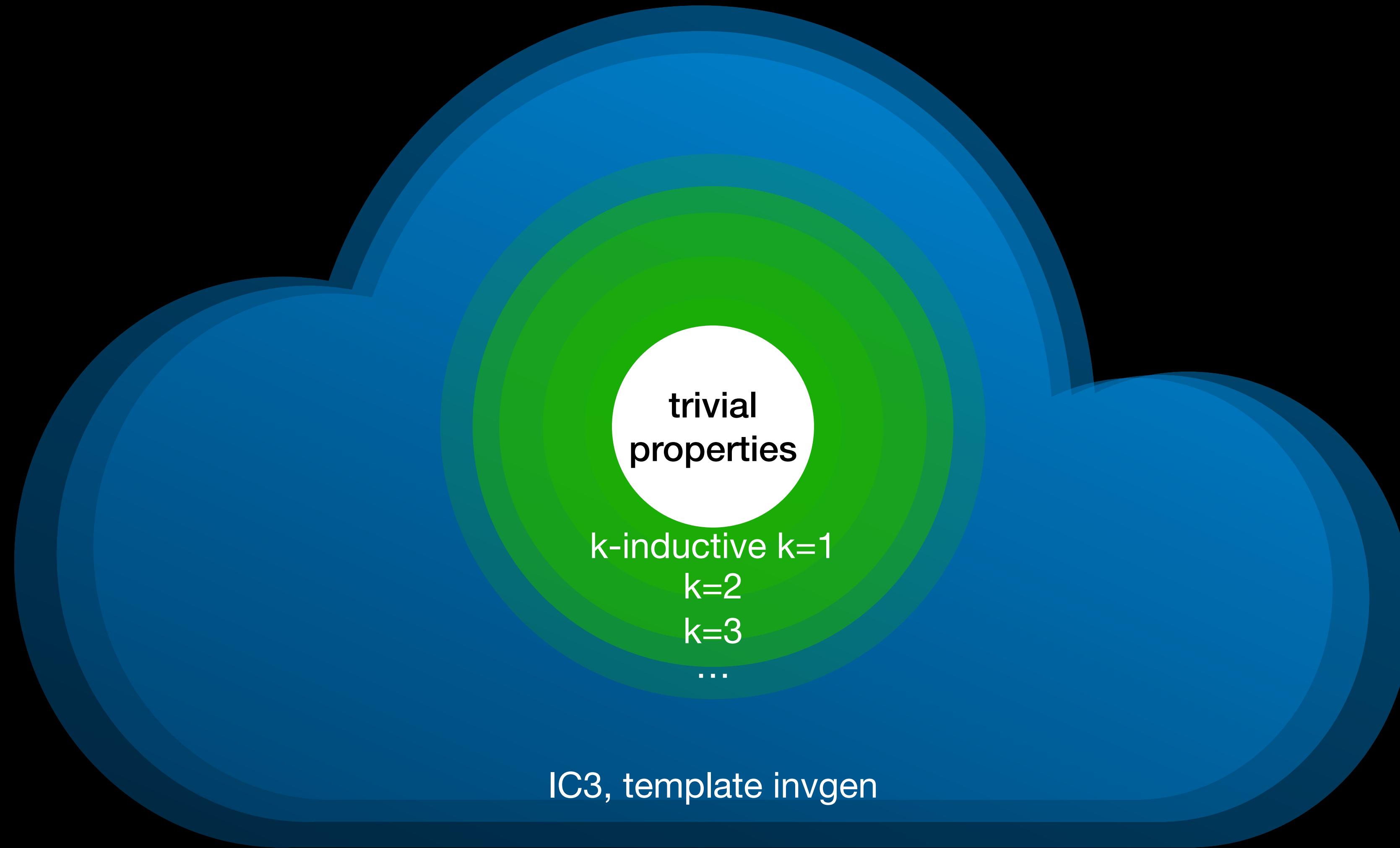
  step(s: state, delta_ok: bool): (state * bool * bool) =
    let
      (state_last_main', last_main) = lastn.step(s.last_main, 10, delta_ok)
      (state_last_signal', last_signal) = lastn.step(s.last_signal, 10, delta_ok)
    in
      ({ init      = false;
         last_main = state_last_main';
         last_signal = state_last_signal' }, last_main, last_signal)
    --%PROPERTY state_last_main' == state_last_signal'
```

# Two problems

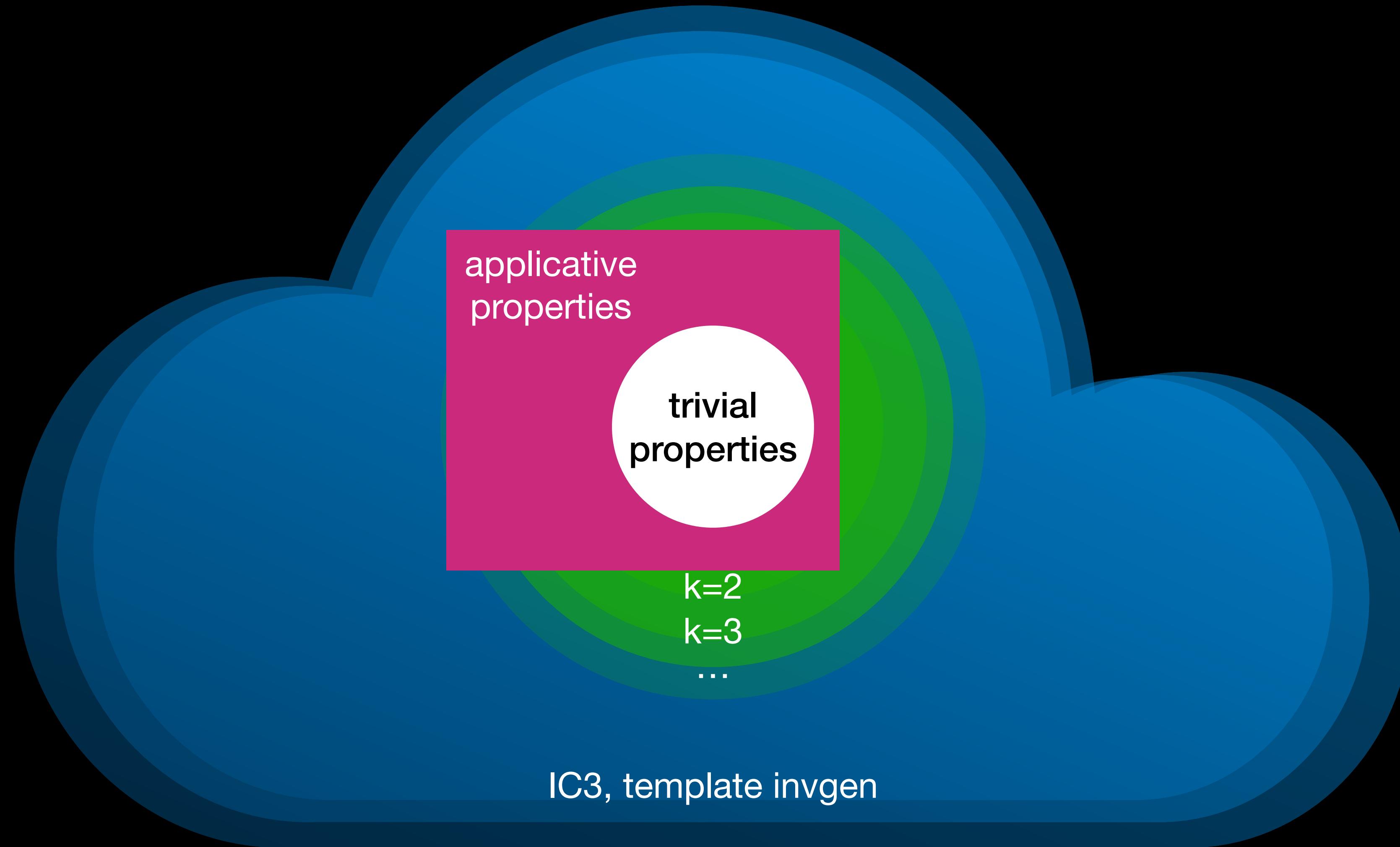
- "obvious" invariants aren't obvious on the translated system
- we can't state the invariant, even though we know it!



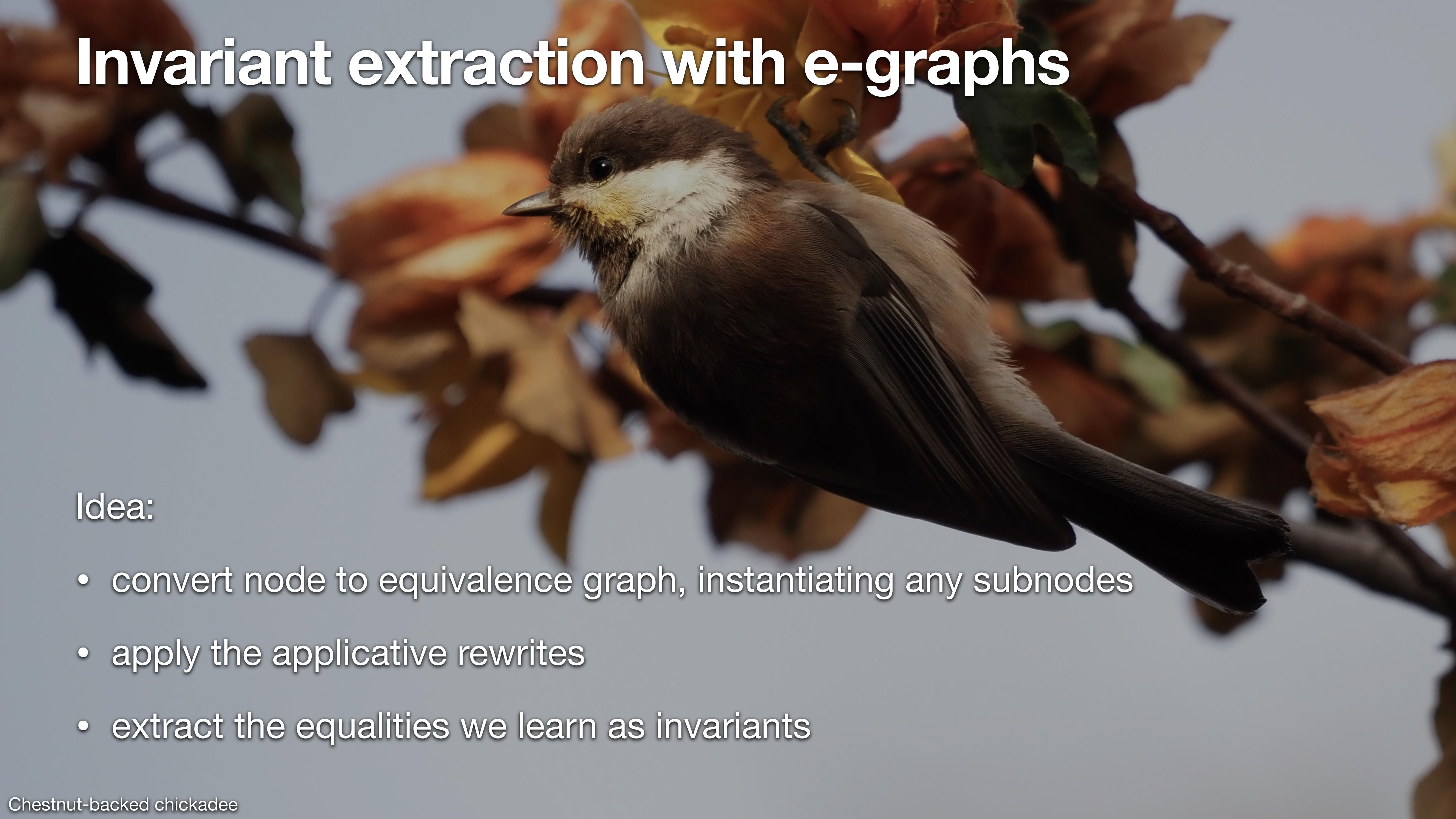
# What we can prove



# What we can prove - proposed class



# Invariant extraction with e-graphs

A close-up photograph of a Chestnut-backed chickadee perched on a branch. The bird has a dark cap, a white forehead, and a chestnut-colored back. It is surrounded by orange flowers and green leaves.

Idea:

- convert node to equivalence graph, instantiating any subnodes
- apply the applicative rewrites
- extract the equalities we learn as invariants

```
node signal_valid(input: int)
(*@contract guarantee not last.out => not ok; *)
    ok
        : bool;
    pre_input
        = pre input;
    del_input
        = 0 -> pre_input;
    subnode last
        = lastn(10, delta_valid(input, del_input))
```

```
node main(sample: SAMPLE)
(*@contract guarantee not last.out => not engaged; *)
    sample_adc
        = sample.adc;
    pre_sample
        = pre sample;
    del_sample
        = SAMPLE_ZERO -> pre_sample;
    del_sample_adc
        = del_sample.adc;
    subnode valid
        = signal_valid(sample_adc);
    subnode last
        = lastn(10, delta_valid(sample_adc, del_sample_adc));
    engaged
        = valid.ok;
```

```
-- subnode valid = signal_valid(sample.adc)
-- guarantee not valid.last.out => not valid.ok;
valid.pre_input = pre sample_adc;
valid.del_input = 0 -> valid.pre_input;
valid.last.out  = lastn(10, delta_valid(sample_adc, valid.del_input));

-- node main(sample)
-- guarantee not last.out => not engaged;
sample_adc      = sample.adc;
pre_sample      = pre sample;

del_sample      = SAMPLE_ZERO -> pre_sample;
del_sample_adc = del_sample.adc;

last.out        = lastn(10, delta_valid(sample_adc, del_sample_adc));
engaged         = valid.ok;
```

```
-- subnode valid = signal_valid(sample.adc)
-- guarantee not valid.last.out => not valid.ok;
valid.pre_input = pre sample_adc;
valid.del_input = 0 -> valid.pre_input;
valid.last.out  = lastn(10, delta_valid(sample_adc, valid.del_input));
```

```
-- node main(sample)
-- guarantee not last.out => not engaged;
sample_adc      = sample.adc;
pre_sample      = pre sample;
```

```
del_sample      = SAMPLE_ZERO -> pre_sample;
del_sample_adc = del_sample.adc;
```

```
last.out        = lastn(10, delta_valid(sample_adc, del_sample_adc));
engaged         = valid.ok;
```

---

```
rewrite prim_distributes_arrow: forall stream e e', pure function f.
f (e -> e') = (f e) -> (f e')
```

```

-- subnode valid = signal_valid(sample.adc)
-- guarantee not valid.last.out => not valid.ok;
valid.pre_input = pre sample_adc;
valid.del_input = 0 -> valid.pre_input;
valid.last.out  = lastn(10, delta_valid(sample_adc, valid.del_input));

-- node main(sample)
-- guarantee not last.out => not engaged;
sample_adc      = sample.adc;
pre_sample      = pre sample;

del_sample      = SAMPLE_ZERO -> pre_sample;
del_sample_adc = del_sample.adc;
del_sample_adc = SAMPLE_ZERO.adc -> pre_sample.adc;

last.out        = lastn(10, delta_valid(sample_adc, del_sample_adc));
engaged         = valid.ok;

```

---

```

rewrite prim_distributes_arrow: forall stream e e', pure function f.
  f (e -> e') = (f e) -> (f e')
with f := _.adc, e := SAMPLE_ZERO, e' := pre sample:
  (SAMPLE_ZERO -> pre sample).adc = SAMPLE_ZERO.adc -> (pre sample).adc

```

```
-- subnode valid = signal_valid(sample.adc)
-- guarantee not valid.last.out => not valid.ok;
valid.pre_input = pre sample_adc;
valid.del_input = 0 -> valid.pre_input;
valid.last.out  = lastn(10, delta_valid(sample_adc, valid.del_input));

-- node main(sample)
-- guarantee not last.out => not engaged;
sample_adc      = sample.adc;
pre_sample      = pre sample;

del_sample      = SAMPLE_ZERO -> pre_sample;
del_sample_adc = del_sample.adc;
del_sample_adc = SAMPLE_ZERO.adc -> pre_sample.adc;

last.out        = lastn(10, delta_valid(sample_adc, del_sample_adc));
engaged         = valid.ok;
```

---

```
rewrite prim_distributes_pre: forall stream e, pure function f.
f (pre e) = pre (f e)
```

```
-- subnode valid = signal_valid(sample.adc)
-- guarantee not valid.last.out => not valid.ok;
valid.pre_input = pre sample_adc;
valid.del_input = 0 -> valid.pre_input;
valid.last.out  = lastn(10, delta_valid(sample.adc, valid.del_input));

-- node main(sample)
-- guarantee not last.out => not engaged;
sample_adc      = sample.adc;
pre_sample      = pre sample;
pre_sample_adc = pre sample_adc;
pre_sample_adc = (pre sample).adc;
del_sample     = SAMPLE_ZERO -> pre_sample;
del_sample_adc = del_sample.adc;
del_sample_adc = SAMPLE_ZERO.adc -> pre_sample_adc;

last.out        = lastn(10, delta_valid(sample.adc, del_sample_adc));
engaged         = valid.ok;
```

---

```
rewrite prim_distributes_pre: forall stream e, pure function f.
  f (pre e) = pre (f e)
with f := _.adc, e := sample
  (pre sample).adc = pre sample.adc
```

```
-- subnode valid = signal_valid(sample.adc)
-- guarantee not valid.last.out => not valid.ok;
valid.pre_input = pre_sample_adc;
valid.del_input = del_sample_adc;
valid.last.out  = lastn(10, delta_valid(sample.adc, del_sample_adc));

-- node main(sample)
-- guarantee not last.out => not engaged;
sample.adc      = sample.adc;
pre_sample      = pre.sample;
pre_sample.adc = pre.sample.adc;
pre_sample.adc = (pre.sample).adc;
del_sample     = SAMPLE_ZERO -> pre_sample;
del_sample.adc = del.sample.adc;
del_sample.adc = SAMPLE_ZERO.adc -> pre_sample.adc;
del_sample.adc = 0 -> pre_sample.adc;
last.out        = lastn(10, delta_valid(sample.adc, del.sample.adc));
engaged         = valid.ok;
```

```
node lastn(const n: int; pred: bool)
pre_count      = pre count;
arr_count      = 0 -> pre_count;
count          = if pred then arr_count + 1 else 0;
out            = count >= n;
```

```
-- subnode last = lastn(10, del_sample_adc)
last.pre_count      = pre last.count;
last.arr_count      = 0 -> last.pre_count;
last.count          = if del_sample_adc
                      then last.arr_count + 1
                      else 0;
last.out            = last.count >= 10;

-- subnode valid.last = lastn(10, del_sample_adc)
valid.last.pre_count = pre valid.last.count;
valid.last.arr_count = 0 -> valid.last.pre_count;
valid.last.count     = if del_sample_adc
                      then valid.last.arr_count + 1
                      else 0;
valid.last.out       = valid.last.count >= 10;
```

```
last.pre_count      = pre last.count;  
last.arr_count      = 0 -> last.pre_count;  
last.count          = if del_sample_adc  
                      then last.arr_count + 1  
                      else 0;
```

==(rewrite recursive binding)=>

```
last.count          = fix (λcount.  
                           if del_sample_adc  
                           then (0 -> pre count) + 1  
                           else 0);
```

```
-- subnode last = lastn(10, delta_valid(...))
last.pre_count      = valid.last.pre_count
                     = pre last.count;
last.arr_count      = valid.last.arr_count
                     = 0 -> last.pre_count;
last.count          = valid.last.count
                     = if del_sample_adc
                         then last.arr_count + 1
                         else 0;
                     = if del_sample_adc
                         then valid.last.arr_count + 1
                         else 0;
                     = fix (λcount.
                           if del_sample_adc
                             then (0 -> pre count) + 1
                             else 0);
last.out            = valid.last.out
                     = last.count >= 10;
```

```
-- subnode last = lastn(10, delta_valid(...))
last.pre_count      = valid.last.pre_count
= pre last.count;
last.arr_count      = valid.last.arr_count
= 0 -> last.pre_count;
last.count          = valid.last.count
= if del_sample_adc
  then last.arr_count + 1
  else 0;
= if del_sample_adc
  then valid.last.arr_count + 1
  else 0;
= fix (λcount.
  if del_sample_adc
  then (0 -> pre count) + 1
  else 0);
last.out            = valid.last.out
= last.count >= 10;
```

invariant:

```
true -> last.pre_count = valid.last.pre_count
```

# Relationship to k-induction



Flannel flower

# Equivalences sometimes stronger

```
node SoFar(pred: bool)
returns (out: bool)
let
  out = pred and (true -> pre out);
tel
```

---

```
assume
  SoFar(X) => P
```

```
show
  SoFar(X) => P
```

- k-induction alone cannot find the invariant that  $\text{SoFar}(X)=\text{SoFar}(X)$

# K-induction sometimes stronger

```
node countmod4()
returns (mod4: int)
let
    mod4 = 0 -> if pre mod4 = 3 then 0 else pre mod4 + 1;
    --%PROPERTY mod4 <= 6;
tel
```

- bound  $\leq 6$  not tight enough to be 1-inductive, but 3-induction can get it

# Modularity

- equivalences can extract invariants from subnodes separately
- could skip equality-saturation on nodes that are too big and still benefit from subnode analyses
- k-induction is monolithic: k parameter controls unfolding for whole system

# Future work

- current implementation in experimental Scala EDSL for Lustre
  - worth implementing in Kind2?
- evaluate on larger programs
- improve clock support
- integration with other invariant generation techniques

